



# Nash bargaining model

$(A, B)$ :  $A, B$   $m \times n$  matrices

## Cooperative Games

Non-transferable utility (NTU)

Players make agreements on what strategies they use.

Sharing of payoffs are not allowed.

Def.

1. Joint strategies:

$$P = \begin{pmatrix} p_{11} & \dots & p_{1n} \\ \vdots & \ddots & \vdots \\ p_{m1} & \dots & p_{mn} \end{pmatrix} \quad m \times n \text{ matrix}$$

①  $P_{ij} \geq 0$

②  $\sum_j P_{ij} = 1$

Probability matrix

$P_{ij}$  = probability that I uses  $i$ -th strategy

and  $\Pi$  uses  $j$ -th strategy

$$\text{Payoff: } (u(P), v(P)) = \left( \sum a_{ij} p_{ij}, \sum b_{ij} p_{ij} \right) \\ = \sum p_{ij} (a_{ij}, b_{ij})$$

In particular, if  $\vec{p} \in \mathcal{P}^m$ ,  $\vec{q} \in \mathcal{P}^n$

$$P = \vec{p}^T \vec{q} = \begin{pmatrix} p_1 q_1 & p_1 q_2 & \dots & p_1 q_n \\ \vdots & \vdots & \ddots & \vdots \\ p_m q_1 & \dots & \dots & p_m q_n \end{pmatrix} \in \mathcal{P}^{m \times n}$$

$m \times n$  matrix

$I$  uses  $\vec{p}$  and  $\Pi$  uses  $\vec{q}$  independently

Not all strategies are of this form e.g.  $\begin{pmatrix} 0.4 & 0 \\ 0 & 0.6 \end{pmatrix}$

2. Cooperative region:  $\mathcal{R} = \{ (u(P), v(P)) : P \in \mathcal{P}^{m \times n} \}$

$$\mathcal{R} = \text{conv} \{ (a_{ij}, b_{ij}) \}$$

① convex polygon in  $\mathbb{R}^2$   
② closed and bounded

3. Status quo point: payoff pair  $(\mu, \nu)$  associated to the solution when  $(A, B)$  is considered as

a non-cooperative game.

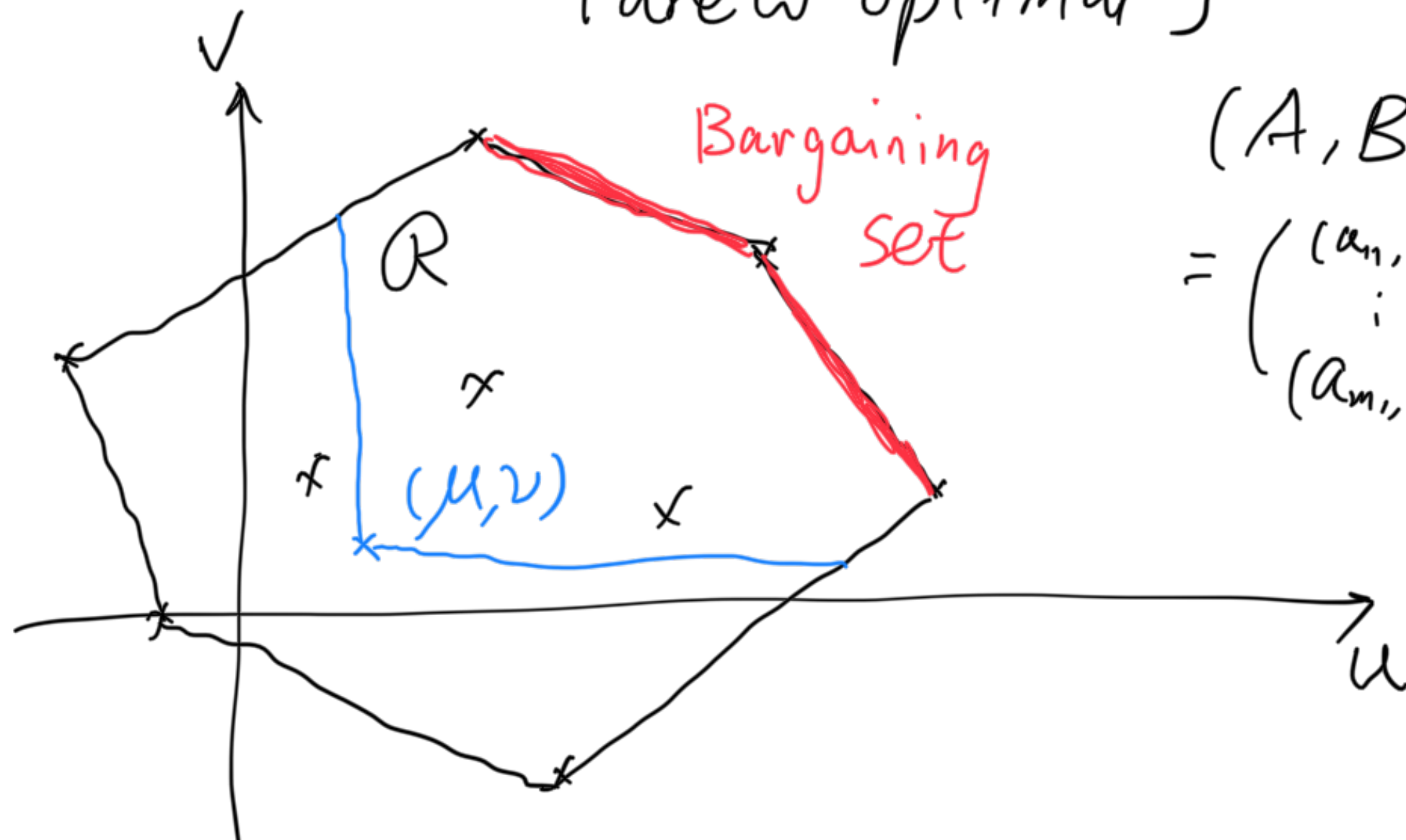
Unless otherwise specified, we take

$$(\mu, \nu) = (v(A), v(B^T)) \text{ security levels}$$

4. We say that  $(u, v) \in \mathcal{R}$  is Pareto optimal if  $u' \geq u, v' \geq v$  and  $(u', v') \in \mathcal{R}$  implies  $(u', v') = (u, v)$

5. Bargaining set:

$$\{(u, v) \in \mathcal{R} : (u, v) \geq (\mu, \nu) \text{ and } (u, v) \text{ is Pareto optimal}\}$$



$$(A, B) = \begin{pmatrix} (a_{11}, b_{11}), \dots, (a_{1n}, b_{1n}) \\ \vdots \\ (a_{m1}, b_{m1}), \dots, (a_{mn}, b_{mn}) \end{pmatrix}$$

# Nash's axiom for bargaining problem

An arbitration is a function  $(\alpha, \beta) = A(\mathcal{Q}, (\mu, \nu))$  defined for a closed and bounded convex set  $\mathcal{Q} \in \mathbb{R}^2$  and  $(\mu, \nu) \in \mathcal{Q}$  such that the following axioms are satisfied.

1. (Individual rationality)  $(\alpha, \beta) \geq (\mu, \nu)$

2. (Pareto optimality) If  $(u, v) \in \mathcal{Q}$  and  $(u, v) \geq (\alpha, \beta)$ , then  $(u, v) = (\alpha, \beta)$ .

3. (Feasibility)  $(\alpha, \beta) \in \mathcal{Q}$

(1), (2), (3)  $\Rightarrow (\alpha, \beta)$  lies on the bargaining set.

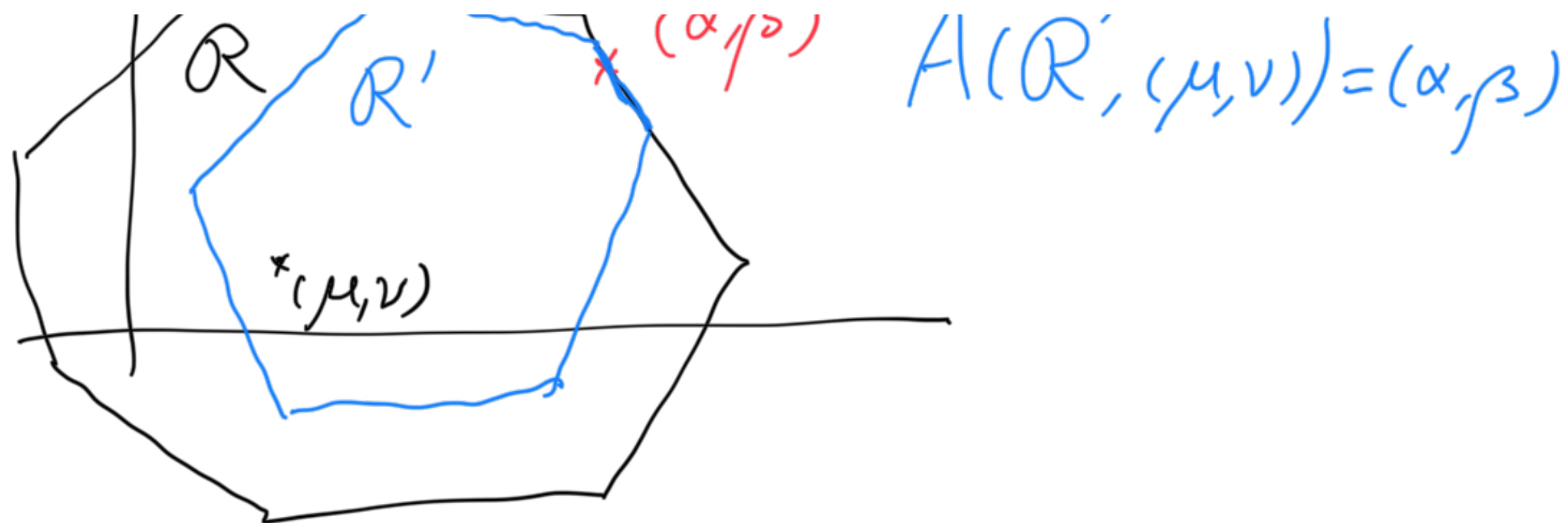
4. (Independence of irrelevant alternatives)

If  $\mathcal{Q}' \subset \mathcal{Q}$  and  $(\mu, \nu), (\alpha, \beta) \in \mathcal{Q}'$ , then

$$A(\mathcal{Q}', (\mu, \nu)) = (\alpha, \beta)$$







5. (Invariant under linear transformations)

$$R' = \{ (u', v') : u' = au + b, v' = cv + d, a, c > 0 \}$$

and  $(u', v') = (au + b, cv + d)$

Then  $A(R', (\mu', \nu')) = (a\alpha + b, c\beta + d)$

6. (Symmetric) If  $R$  is symmetric, i.e.,  
 $(u, v) \in R \Rightarrow (v, u) \in R$ , and  $\mu = \nu$ , then

$$\alpha = \beta$$



Example:

1. Dating game:  $(A, B) = \begin{pmatrix} (4, 2) & (0, 0) \\ (0, 0) & (1, 3) \end{pmatrix}$

$$A = \begin{pmatrix} 4 & 0 \\ 0 & 1 \end{pmatrix} \begin{matrix} 4 \\ -1 \end{matrix} \times \begin{matrix} 0.25 \\ 0.75 \end{matrix}$$

$$\begin{matrix} 4 & -1 \\ 0.25 & 0.75 \end{matrix} \times$$

$$B = \begin{pmatrix} 2 & 0 \\ 0 & 3 \end{pmatrix}$$

$$\textcircled{2} \left( \frac{18}{5}, \frac{32}{15} \right) = p(4, 2) + (1-p)(1, 3)$$

$$v = v(B^T) = 1.2$$

$$\mu = v(A) = 0.8$$

$$(\mu, v) = (0.8, 1.2)$$

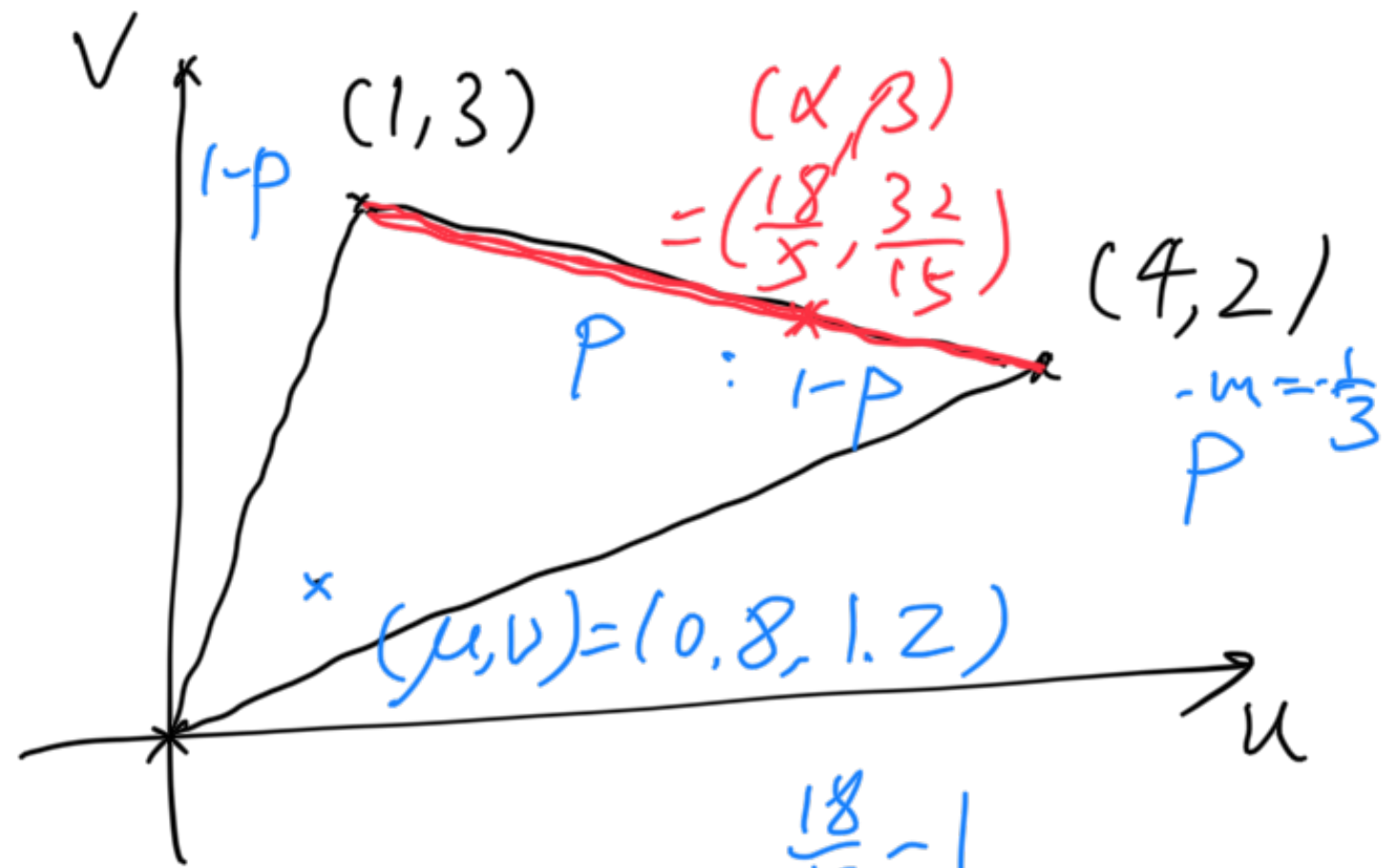
$$v - 3 = -\frac{1}{3}(u - 1)$$

$$v = -\frac{1}{3}u + \frac{10}{3}$$

$$g(u, v) = (u - \mu)(v - v)$$

$$= (u - 0.8)(v - 1.2)$$

$$= \left( u - \frac{4}{5} \right) \left( -\frac{1}{3}u + \frac{10}{3} - \frac{6}{5} \right)$$



$\textcircled{1}$

$$p = \frac{\frac{18}{5} - 1}{4 - 1} = \frac{3 - \frac{32}{15}}{3} = \frac{12}{15} = \frac{4}{5}$$

$$= -\frac{1}{3}u^2 + \frac{12}{5}u - \frac{128}{75} \quad \frac{3-2}{15} = \frac{1}{15}$$

$(\alpha, \beta)$  is the point on the bargaining where  $g(u, v) = (u - \mu)(v - \nu)$  attains its maximum.

$$\alpha = \frac{\frac{12}{5}}{2 \times \frac{1}{3}} = \frac{18}{5}, \quad \beta = -\frac{1}{3} \times \frac{18}{5} + \frac{10}{3} = \frac{32}{15}$$

$\therefore$  The arbitration pair is  $(\alpha, \beta) = \left(\frac{18}{5}, \frac{32}{15}\right)$

The joint strategy is

$$P = \begin{pmatrix} p & 0 \\ 0 & 1-p \end{pmatrix} \quad \text{where} \quad p = \frac{\frac{18}{5} - 1}{4 - 1} = \frac{13}{15}$$

$$= \begin{pmatrix} \frac{13}{15} & 0 \\ 0 & \frac{2}{15} \end{pmatrix}$$

$$2. \quad (A, B) = \begin{pmatrix} (2, 3) & (1, 5) & (6, 1) \\ (3, 4) & (1, 2) & (0, 3) \end{pmatrix}$$



$$(\mu, \nu) = \left(1, \frac{7}{2}\right) = (\nu(A), \nu(B^T))$$

$$\textcircled{1} \quad \nu - 4 = -(u - 3)$$

$$\nu = -u + 7$$

$$g(u, \nu)$$

$$= (u - 1) \left(\nu - \frac{7}{2}\right)$$

$$= (u - 1) \left(-u + 7 - \frac{7}{2}\right)$$

$$= (u - 1) \left(-u + \frac{7}{2}\right)$$

$$= -u^2 + \frac{9}{2}u - \frac{7}{2}$$

maximum when  $u = \frac{+\frac{9}{2}}{2} = +\frac{9}{4} = 2.25 < 3$

$\left(\frac{9}{4}, \frac{19}{4}\right)$  lies outside  $\mathcal{R}$

$$\begin{cases} \nu - 4 = -(u - 3) \\ \nu - \frac{7}{2} = (u - 1) \end{cases}$$

$$\Rightarrow (u, \nu) = \left(\frac{9}{4}, \frac{19}{4}\right)$$

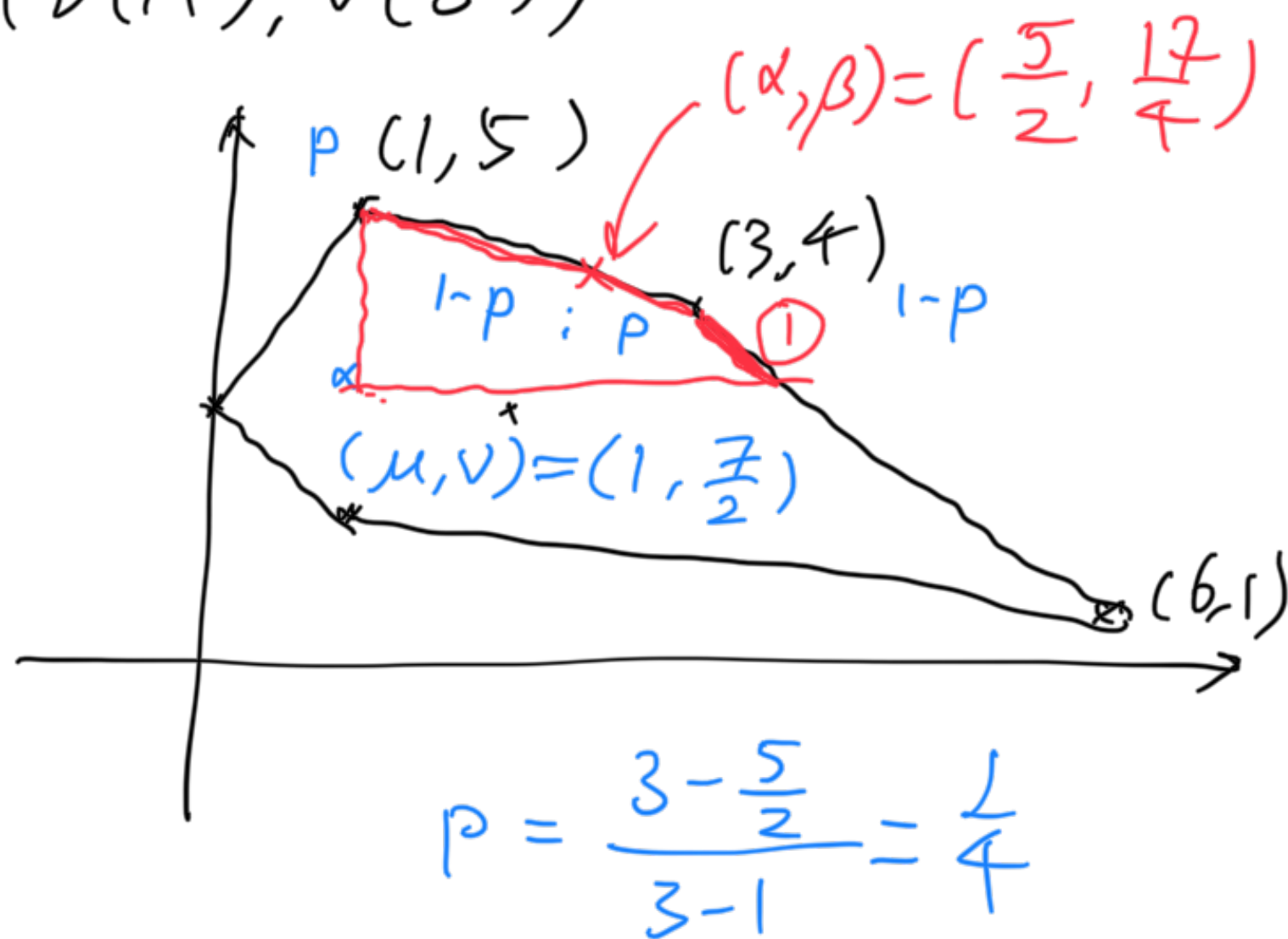
$\textcircled{2}$

$$\nu - 5 = -\frac{1}{2}(u - 1)$$

$$\nu = -\frac{1}{2}u + \frac{11}{2}$$

$$g(u, \nu) = (u - 1) \left(\nu - \frac{7}{2}\right) = (u - 1) \left(-\frac{1}{2}u + 2\right)$$

$$= -\frac{1}{2}u^2 + \frac{5}{2}u - 2$$



maximum when  $u = \frac{5}{2}$

$(\frac{5}{2}, \frac{17}{4})$  lies on the line segment.

$\therefore$  The arbitration pair is  $(\alpha, \beta) = (\frac{5}{2}, \frac{17}{4})$

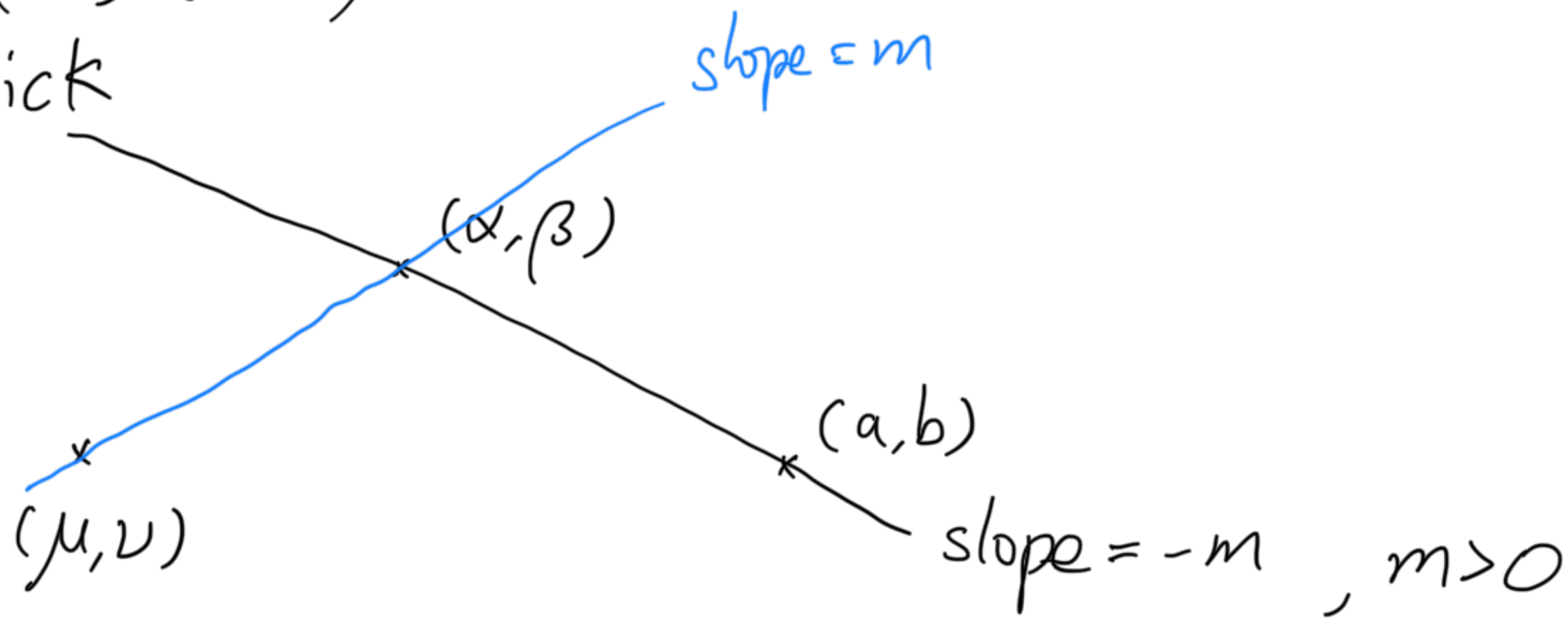
The joint strategy:

$$P = \begin{pmatrix} 0 & p & 0 \\ 1-p & 0 & 0 \end{pmatrix}$$

where  $p = \frac{3 - \frac{5}{2}}{3-1} = \frac{1}{4}$

$$= \begin{pmatrix} 0 & 0.25 & 0 \\ 0.75 & 0 & 0 \end{pmatrix}$$

Little trick



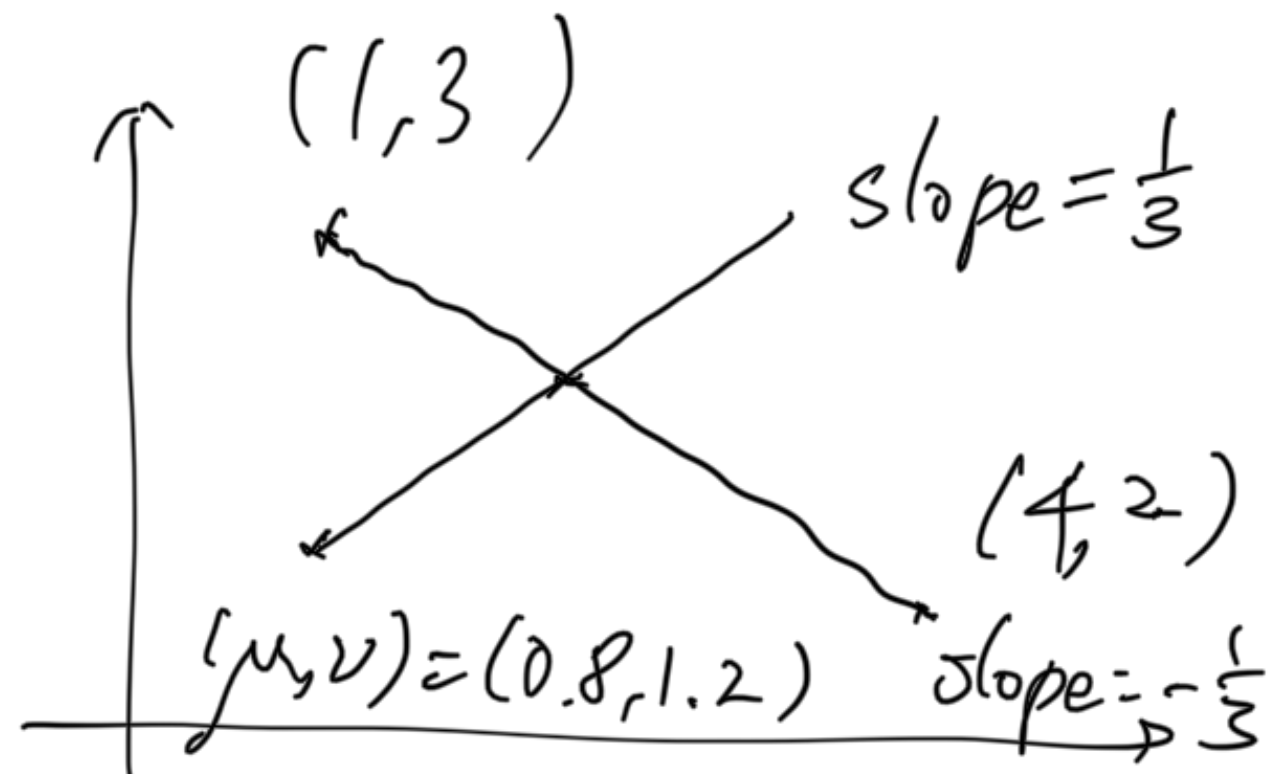
$$\beta - b = -m(\alpha - a)$$

$$\beta - v = m(\alpha - \mu)$$

In dating game:

$$\begin{cases} v - 3 = -\frac{1}{3}(u - 1) \\ v - 1.2 = \frac{1}{3}(u - 0.8) \end{cases}$$

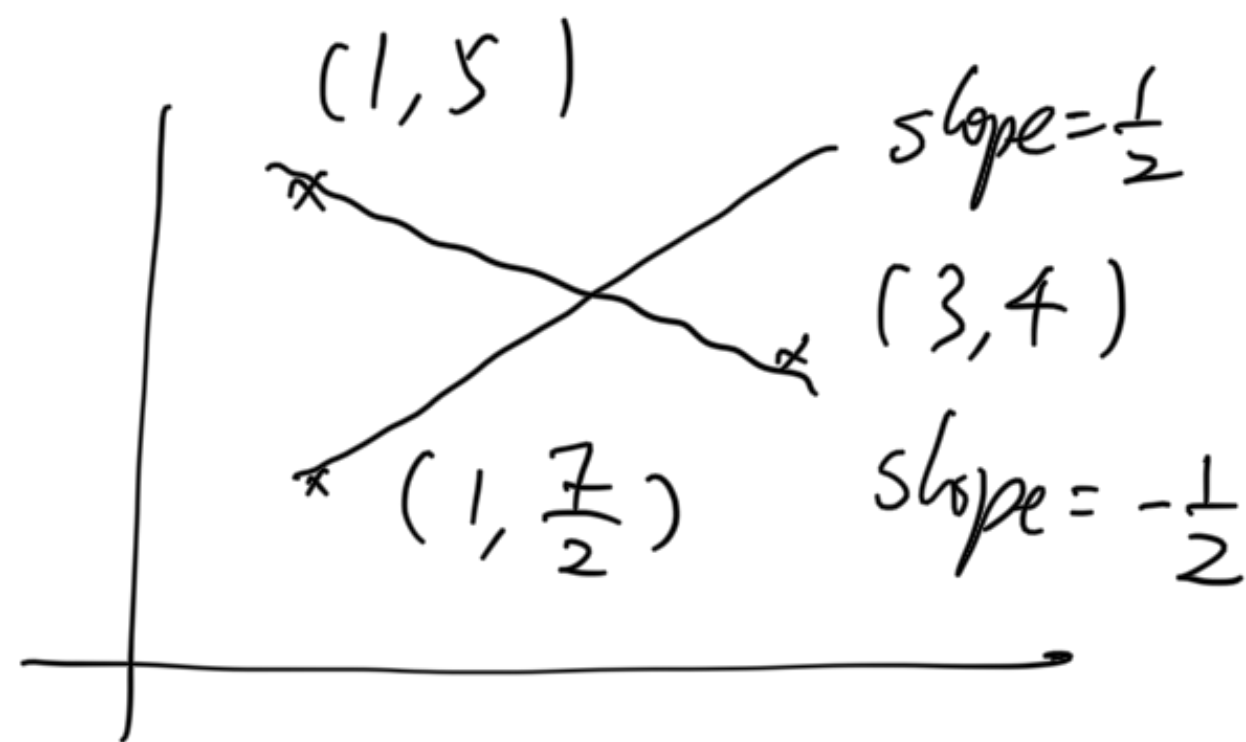
$$\Rightarrow (u, v) = \left(\frac{18}{5}, \frac{32}{15}\right)$$



In the second example

$$\begin{cases} v - 5 = -\frac{1}{2}(u - 1) \\ v - \frac{7}{2} = \frac{1}{2}(u - 1) \end{cases}$$

$$\Rightarrow (u, v) = \left(\frac{5}{2}, \frac{17}{4}\right)$$



Thm (Nash) There exists a unique arbitration

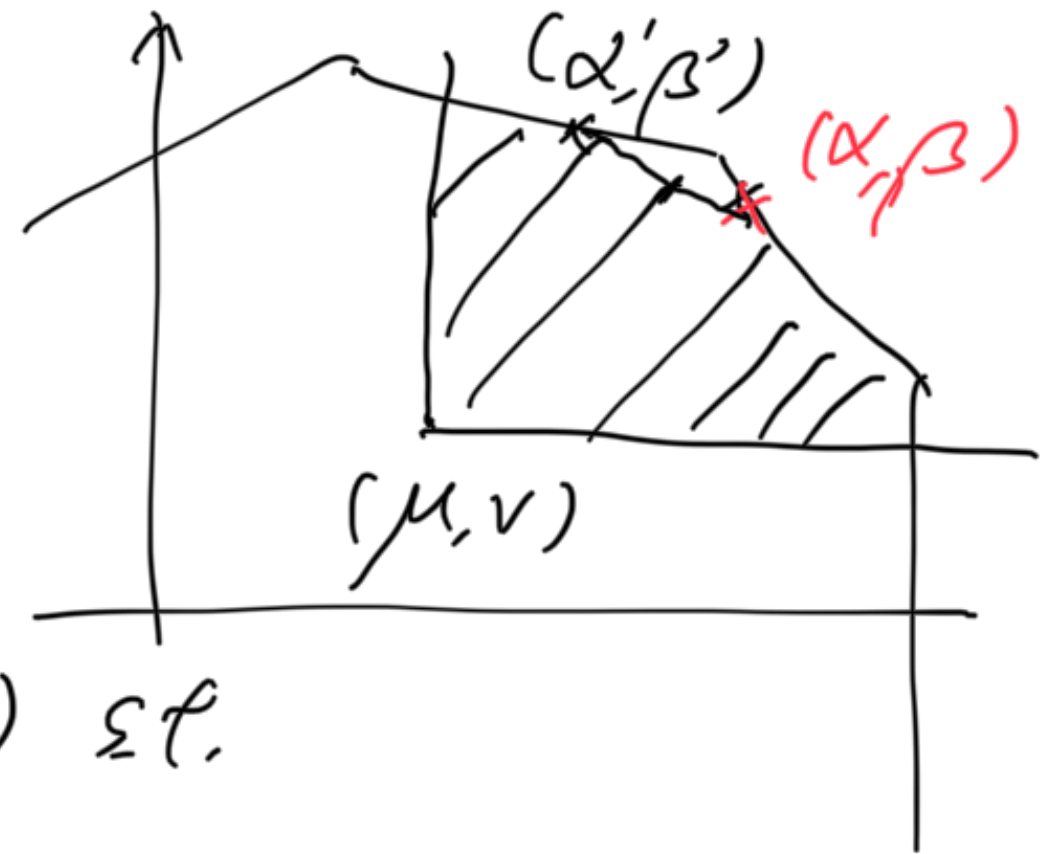
$A(\mathcal{R}, (\mu, v))$  satisfies the Nash's axioms.

Lemma: Define

$$g(u, v) = (u - \mu)(v - \nu) \text{ for } (u, v) \in \mathbb{R} \text{ with } (u, v) \succeq (\mu, \nu)$$

Then  $\exists$  unique  $(\alpha, \beta)$  s.t.

$$g(\alpha, \beta) = \max_{\substack{(u, v) \in \mathbb{R} \\ (u, v) \succeq (\mu, \nu)}} g(u, v)$$



Proof. Suppose  $\exists (\alpha', \beta') \neq (\alpha, \beta)$  s.t.

$$g(\alpha', \beta') = g(\alpha, \beta)$$

Then either  $\alpha' > \alpha$  and  $\beta' < \beta$  or  $\alpha' < \alpha$  and  $\beta' > \beta$

$$\begin{aligned} & g\left(\frac{\alpha + \alpha'}{2}, \frac{\beta + \beta'}{2}\right) \\ &= \left(\frac{\alpha + \alpha'}{2} - \mu\right)\left(\frac{\beta + \beta'}{2} - \nu\right) \\ &= \frac{1}{4} \left( (\alpha - \mu) + (\alpha' - \mu) \right) \left( (\beta - \nu) + (\beta' - \nu) \right) \\ &= \frac{1}{4} \left[ (\alpha - \mu)(\beta - \nu) + (\alpha - \mu)(\beta' - \nu) + (\alpha' - \mu)(\beta - \nu) + (\alpha' - \mu)(\beta' - \nu) \right] \end{aligned}$$



$$= \frac{1}{4} \left[ (\alpha - \mu)(\beta - \nu) + (\alpha - \mu)(\beta - \nu + \beta' - \beta) + (\alpha' - \mu)(\beta' - \nu + \beta - \beta') + (\alpha' - \mu)(\beta' - \nu) \right]$$

$$= \frac{1}{4} \left[ g(\alpha, \beta) + \underline{(\alpha - \mu)(\beta - \nu)} + (\alpha - \mu)(\beta' - \beta) + \underline{(\alpha' - \mu)(\beta' - \nu)} + (\alpha' - \mu)(\beta - \beta') + g(\alpha', \beta') \right]$$

$$= \frac{1}{4} \left[ 2g(\alpha, \beta) + 2g(\alpha', \beta') + (\beta' - \beta)(\alpha - \mu) - (\alpha' - \mu)(\beta - \beta') \right]$$

$$= \frac{1}{4} \left[ 2g(\alpha, \beta) + 2g(\alpha', \beta') + (\beta' - \beta)(\alpha - \alpha') \right]$$

$$> g(\alpha, \beta)$$

Contradicts  $\left(\frac{\alpha + \alpha'}{2}, \frac{\beta + \beta'}{2}\right) \in \mathcal{R}$

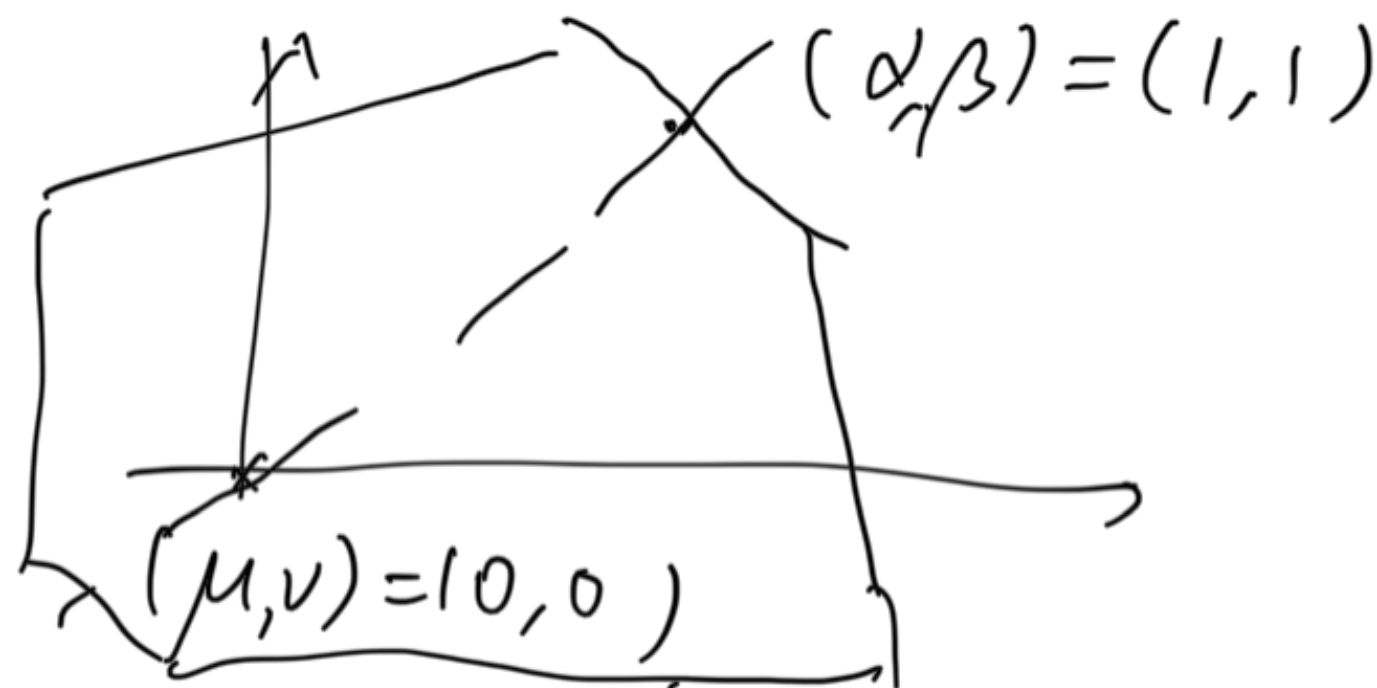
and  $g$  attains its maximum at  $(\alpha, \beta)$ .

$\square$   $\mathcal{R} \cap (\mathcal{R} \cup \dots) = (\alpha, \beta)$  at

Define  $A(\mu, \nu) = (\alpha, \beta)$  s.t.

$$g(\alpha, \beta) = \max_{\substack{(u, v) \in \mathcal{R} \\ (u, v) \geq (\mu, \nu)}} g(u, v)$$

To prove uniqueness, suppose  $A'$  is another arbitration function. By taking a linear transformation, we may assume  $(\mu, \nu) = (0, 0)$  and  $(\alpha, \beta) = (1, 1)$



Use 4, 6  $\Rightarrow A'(\mathcal{R}, (0, 0)) = (1, 1)$